

Analysis of Flexible Serial Lines with Setups

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June 2, 2015

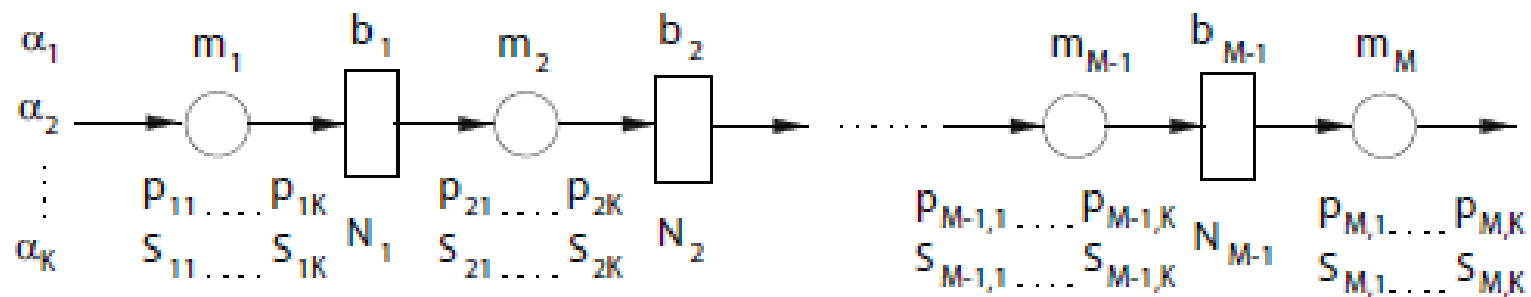


Outline

- System Modeling
- Single Machine Line
- Two-Machine Line
- Long Serial Line
- Bottleneck Analysis
- Case study
- Summary

System Modeling

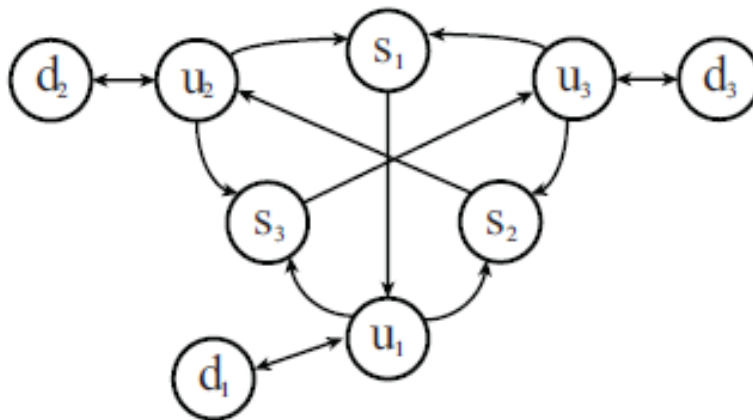
- Serial Production Line with M machines and K products



- Incoming sequence
 - product type: discrete distribution with probability α_j to be type j
 - $\sum_{j=1}^K \alpha_j = 1$
- Bernoulli reliability model
 - Parameters p_{ij} , for machine i product j
- Bernoulli setup model
 - Setup success with probability s_{ij}

Single Machine

- A single machine with $K = 3$ products.
- State transition highlights
 - Failed machine must return to the same type up state.
 - Product changeover must go through setup.
 - Self loops are ignored.



Transition diagram

$$\pi_{u_{1j}} = \frac{\alpha_j}{1 + \sum_{l=1}^K \frac{\alpha_l(1-\alpha_l)}{s_{1l}} + \alpha_l^2 \left(\frac{1}{p_{1l}} - 1 \right)}$$

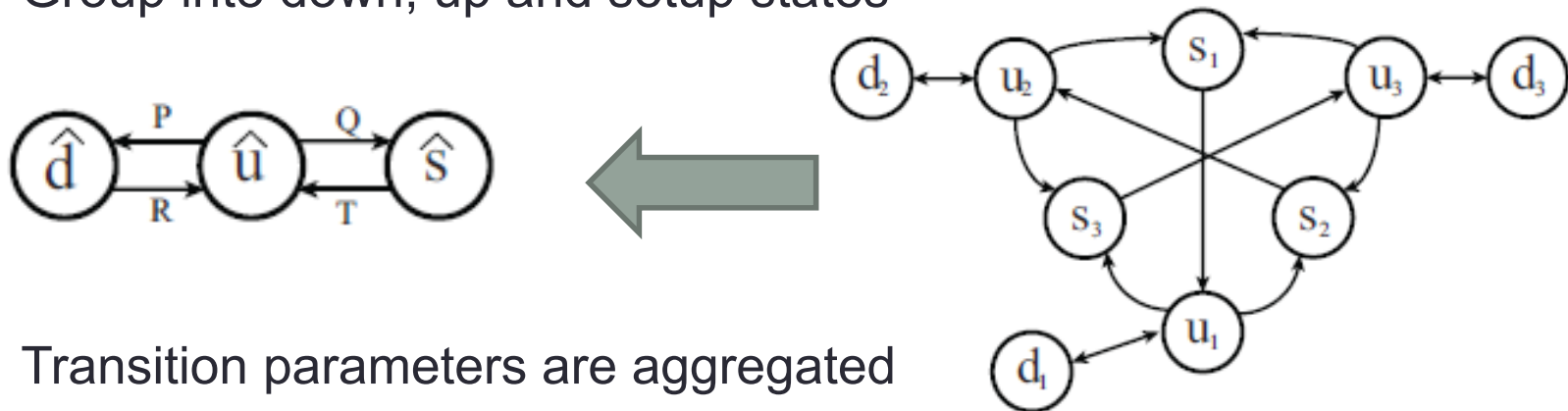
$$\pi_{s_{1j}} = \frac{s_{1j}}{1 + \sum_{l=1}^K \frac{\alpha_l(1-\alpha_l)}{s_{1l}} + \alpha_l^2 \left(\frac{1}{p_{1l}} - 1 \right)}$$

$$\pi_{d_{1j}} = \frac{\alpha_j^2 \left(\frac{1}{p_{1j}} - 1 \right)}{1 + \sum_{l=1}^K \frac{\alpha_l(1-\alpha_l)}{s_{1l}} + \alpha_l^2 \left(\frac{1}{p_{1l}} - 1 \right)}$$

Steady state distribution

Single Machine State Aggregation

- Scale up the analytical method for long line
 - Group into down, up and setup states



- Transition parameters are aggregated

$$P = \frac{\sum_j (1 - p_{1j}) \alpha_j \pi_{u_{1j}}}{\sum_j \pi_{u_{1j}}}$$

Breaks down (pointing to $(1 - p_{1j})$)

Still make product j (pointing to α_j)

$$R = \frac{\sum_j p_{1j} \pi_{d_{1j}}}{\sum_j \pi_{d_{1j}}}$$

$$Q = \frac{\sum_j (1 - \alpha_j) \pi_{u_{1j}}}{\sum_j \pi_{u_{1j}}}$$

From an up state (pointing to $\pi_{u_{1j}}$)
$$T = \frac{\sum_j s_{ij} \pi_{s_{1j}}}{\sum_j \pi_{s_{1j}}}$$

Single Machine State Aggregation

- Parameter aggregation

$$P = \frac{\sum_j (1 - p_{1j}) \alpha_j \pi_{u_{1j}}}{\sum_j \pi_{u_{1j}}} = \sum_j (1 - p_{1j}) \alpha_j^2,$$

$$R = \frac{\sum_j p_{1j} \pi_{d_{1j}}}{\sum_j \pi_{d_{1j}}} = \frac{P}{\sum_j \alpha_j^2 (1/p_{1j} - 1)},$$

$$Q = \frac{\sum_j (1 - \alpha_j) \pi_{u_{1j}}}{\sum_j \pi_{u_{1j}}} = \sum_j (1 - \alpha_j) \alpha_j,$$

$$T = \frac{\sum_j s_{ij} \pi_{s_{1j}}}{\sum_j \pi_{s_{1j}}} = \frac{Q}{\sum_j (1 - \alpha_j) \alpha_j / s_{ij}}.$$

- Steady states are computed and matched with original machine.

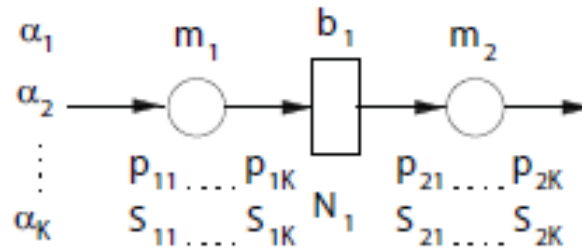
$$\pi_{\hat{u}} = \frac{1}{P/R + 1 + Q/T}, \quad \pi_{\hat{s}} = \frac{Q/T}{P/R + 1 + Q/T}, \quad \pi_{\hat{d}} = \frac{P/R}{P/R + 1 + Q/T}.$$

- Throughput for product type j

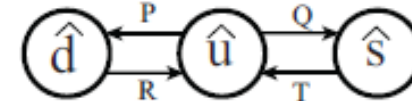
- $\pi_{u_{1j}} = \alpha_j \pi_{\hat{u}}$

Two-Machine Line

- Full state space size: $9K^2(N^K + 1)$
- Aggregated state space size: $9(N + 1)$



For each of the machines



- Performance evaluation (subscripts P, R, Q, T are machine indexes for m_1 and m_2):

$$PR = \Phi_p(P_1, P_2, R_1, R_2, Q_1, Q_2, T_1, T_2, N) = \pi_{N,uu} + \sum_{k=0}^{N-1} (\pi_{k,us} + \pi_{k,ud} + \pi_{k,uu})$$

$$= \sum_{k=1}^N (\pi_{k,su} + \pi_{k,du} + \pi_{k,uu}),$$

$$BL = \Phi_b(P_1, P_2, R_1, R_2, Q_1, Q_2, T_1, T_2, N) = \pi_{N,us} + \pi_{N,ud},$$

$$ST = \Phi_s(P_1, P_2, R_1, R_2, Q_1, Q_2, T_1, T_2, N) = \pi_{0,uu} + \pi_{0,su} + \pi_{0,du},$$

$$WIP = \Phi_w(P_1, P_2, R_1, R_2, Q_1, Q_2, T_1, T_2, N) = \sum_k \sum_{m_1} \sum_{m_2} k \cdot \pi_{k,m_1 m_2}.$$

Long Serial Line

- Use blockage and starvation information to adjust production failure parameters P_i and R_i .

Procedure 1

$$\begin{aligned}
 bl_i^b(n+1) &= \Phi_b(P_i^f(n), P_{i+1}^b(n+1), R_i^f(n), R_{i+1}^b(n+1), Q_i, Q_{i+1}, T_i, T_{i+1}, N_i), \\
 st_i^f(n+1) &= \Phi_s(P_{i-1}^f(n), P_i^b(n+1), R_{i-1}^f(n), R_i^b(n+1), Q_{i-1}, Q_i, T_{i-1}, T_i, N_{i-1}), \\
 P_i^b(n+1) &= P_i + R_i bl_i^b(n+1), \quad R_i^b(n+1) = R_i(1 - bl_i^b(n+1)), \quad i = 1, \dots, M, \\
 P_i^f(n+1) &= P_i + R_i st_i^f(n+1), \quad R_i^f(n+1) = R_i(1 - st_i^f(n+1)), \quad i = 1, \dots, M
 \end{aligned}$$

with initial conditions

$$P_i^f(0) = P_i, \quad R_i^f(0) = R_i, \quad i = 1, \dots, M$$

and boundary conditions

$$P_1^f(n) = P_1, \quad P_M^b(n) = P_M, \quad R_1^f(n) = R_1, \quad R_M^b(n) = R_M, \quad n = 0, 1, 2, \dots$$

- The algorithm converges.

Accuracy Evaluation

- 10,000 production lines are randomly generated.

$$\begin{aligned}
 M &\in \{2, \dots, 10\}, & K &\in \{2, \dots, 5\}, \\
 p_{ij} &\in (0.7, 0.99), & i &= 1, \dots, M, j = 1, \dots, K, \\
 s_{ij} &\in (0.5, 0.99), & i &= 1, \dots, M, j = 1, \dots, K, \\
 N_i &\in \{5, 6, \dots, 15\}, & i &= 1, \dots, M - 1, \\
 \alpha_j &\in (0, 1), & j &= 1, \dots, K, \quad \text{s.t.} \sum_{j=1}^K \alpha_j = 1.
 \end{aligned}$$

- Computation results are compared with simulation results.
 - Production rate: avg err 2.90%
 - WIP: avg err 9.40%
 - BL,ST: avg err 0.02

$$\begin{aligned}
 \delta_{PR} &= \frac{|PR^{sim} - PR^{model}|}{PR^{sim}} \cdot 100\%, \\
 \delta_{WIP_i} &= \frac{|WIP_i^{sim} - WIP_i^{model}|}{N_i} \cdot 100\%, \\
 \delta_{BL_i} &= |BL_i^{sim} - BL_i^{model}|, \\
 \delta_{ST_i} &= |ST_i^{sim} - ST_i^{model}|,
 \end{aligned}$$

Bottleneck Product Type in Setup Time

- Focus on one machine, find the most severe setup bottleneck type.

DEFINITION 1. Under assumption 1-9, product type j on machine m_i is the setup bottleneck product type if

$$\left| \frac{\partial PR}{\partial T_{s_{ij}}} \right| > \left| \frac{\partial PR}{\partial T_{s_{il}}} \right|, \quad \forall l \neq j, l, j \in \{1, \dots, K\},$$

where $T_{s_{ij}} = \frac{1}{s_{ij}}$ is the setup time.

- Recall

$$Q = \frac{\sum_j (1 - \alpha_j) \pi_{u_{1j}}}{\sum_j \pi_{u_{1j}}} = \sum_j (1 - \alpha_j) \alpha_j,$$

$$T = \frac{\sum_j s_{ij} \pi_{s_{1j}}}{\sum_j \pi_{s_{1j}}} = \frac{Q}{\sum_j (1 - \alpha_j) \alpha_j / s_{ij}}$$

$$\pi_{\hat{u}} = \frac{1}{P/R + 1 + Q/T}$$

- Indicator

PROPOSITION 3. Under assumptions 1-9, for machine m_i , the setup bottleneck product type j is the one with the largest value of $(1 - \alpha_j) \alpha_j$, or equivalently, the smallest value of $|\alpha_j - 0.5|$.

Selectivity of Setup and Failure Improvement

- On a given machine
 - focus on setup time reduction or downtime reduction?

DEFINITION 2. Under assumptions 1-9, for machine m_i , setup time reduction has higher selectivity in continuous improvement if

$$\sum_{j=1}^K \left| \frac{\partial PR}{\partial T_{s_{ij}}} \right| > \sum_{j=1}^K \left| \frac{\partial PR}{\partial T_{d_{ij}}} \right|, \quad \forall i \in \{1, \dots, M\}.$$

- Recall

$$\pi_{u_{1j}} = \frac{\alpha_j}{1 + \sum_{l=1}^K \frac{\alpha_l(1-\alpha_l)}{s_{1l}} + \alpha_l^2 \left(\frac{1}{p_{1l}} - 1 \right)}$$

- Indicator

PROPOSITION 4. Under assumptions 1-9 with $M = 1$, setup time reduction has higher selectivity if and only if $\sum_j \alpha_j^2 < 0.5$.

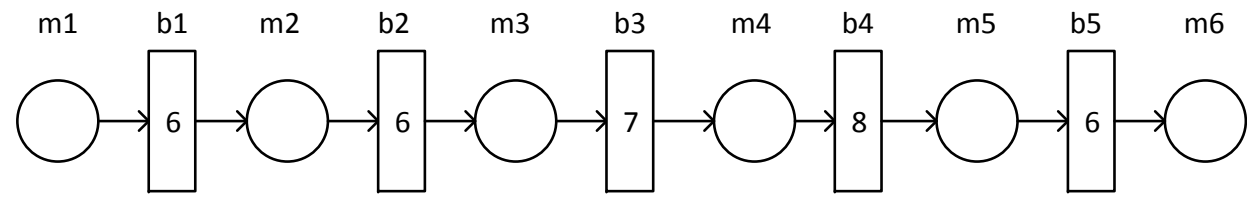
- Accuracy for long line ($M \geq 2$)
 - 84.54% out of 10,000 cases, it finds the right selectivity.
 - In the failed cases, derivatives are small.
 - 92.64%: all machines in a product line have the same selectivity.

Showcase and Effectiveness

p_{i1}	0.9708	0.9834	0.8473	0.8468	0.8013	0.9700
p_{i2}	0.8108	0.7334	0.9341	0.8169	0.7725	0.8212
p_{i3}	0.7289	0.7396	0.9826	0.9868	0.8726	0.7179
S_{i1}	0.6939	0.7413	0.9285	0.6062	0.6172	0.6676
S_{i2}	0.8596	0.8927	0.8591	0.7804	0.8188	0.7185
S_{i3}	0.8979	0.6756	0.8747	0.6734	0.7474	0.8502

α_3 most close to 0.5, product 3 is bottleneck

α_1	0.4055
α_2	0.1303
α_3	0.4642



$\Delta PR / \Delta T_{s_{i1}}$	0.0000	0.0001	0.0003	0.0207	0.0666	0.0185
$\Delta PR / \Delta T_{s_{i2}}$	0.0000	0.0000	0.0001	0.0098	0.0314	0.0086
$\Delta PR / \Delta T_{s_{i3}}$	0.0000	0.0001	0.0003	0.0214	0.0688	0.0191
$\sum_j \Delta PR / \Delta T_{s_{ij}}$	0.0000	0.0001	0.0008	0.0519	0.1667	0.0462
$\sum_j \Delta PR / \Delta T_{d_{ij}}$	0.0000	0.0001	0.0003	0.0285	0.1067	0.0317

$\sum \alpha_j^2 < 0.5$
Setup time reduction is more effective

Bottleneck Setup Machine

- On a serial line
 - Which machine is most impeding in terms of setup time?

DEFINITION 3. Machine m_i is the joint setup bottleneck machine if

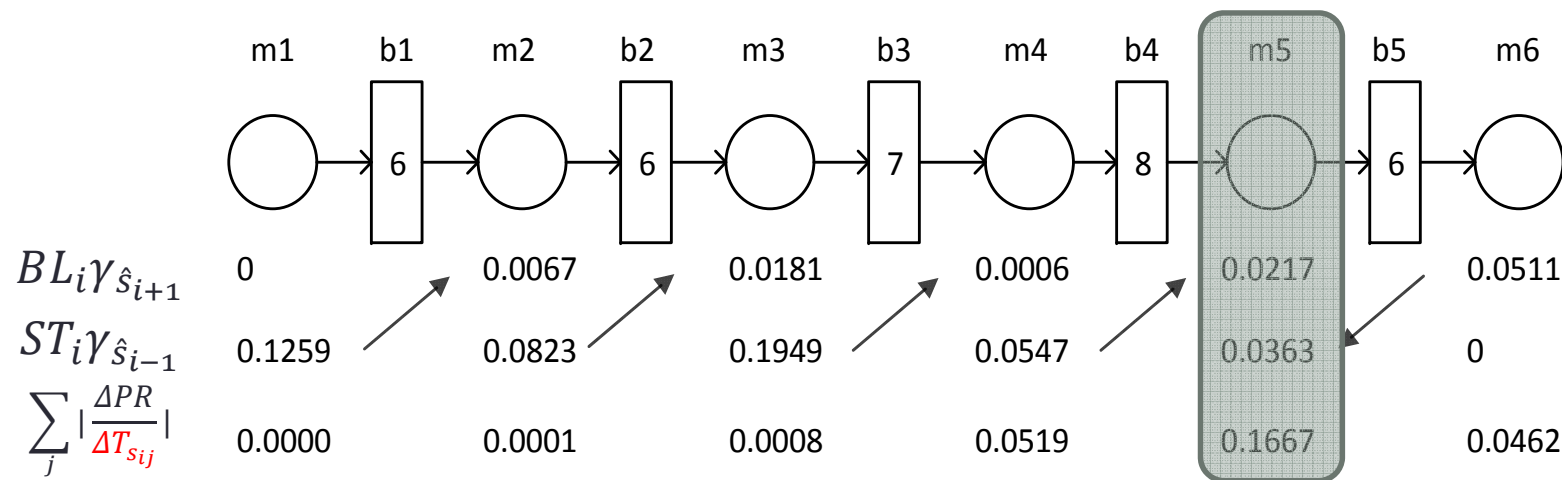
$$\sum_{j=1}^K \left| \frac{\partial PR}{\partial T_{s_{ij}}} \right| > \sum_{j=1}^K \left| \frac{\partial PR}{\partial T_{s_{lj}}} \right|, \quad \forall i \neq l \in \{1, \dots, M\}.$$

- Indicator is based on the adjusted starvation and blockage.
 - Use downstream setup percentage ($\gamma^{\hat{s}_{i+1}}$) in idle time to adjust blockage/starvation
 - Use upstream setup percentage ($\gamma^{\hat{s}_{i-1}}$) in idle time to adjust blockage/starvation
- Arrow assignment rule
 - $BL_i \gamma^{\hat{s}_{i+1}} > ST_{i+1} \gamma^{\hat{s}_i}$, assign the arrow from m_i to m_{i+1} .
 - $BL_i \gamma^{\hat{s}_{i+1}} < ST_{i+1} \gamma^{\hat{s}_i}$, assign the arrow from m_{i+1} to m_i .
- In case of multiple bottleneck, severity index is used.

$$\begin{aligned} S_i &= |ST_{i+1} \gamma^{\hat{s}_i} - BL_i \gamma^{\hat{s}_{i+1}}| + |BL_{i-1} \gamma^{\hat{s}_i} - ST_i \gamma^{\hat{s}_{i-1}}|, \quad i = 2, \dots, M-1, \\ S_1 &= |ST_2 \gamma^{\hat{s}_1} - BL_1 \gamma^{\hat{s}_2}|, \\ S_M &= |BL_{M-1} \gamma^{\hat{s}_M} - ST_M \gamma^{\hat{s}_{M-1}}|. \end{aligned}$$

Showcase and Effectiveness

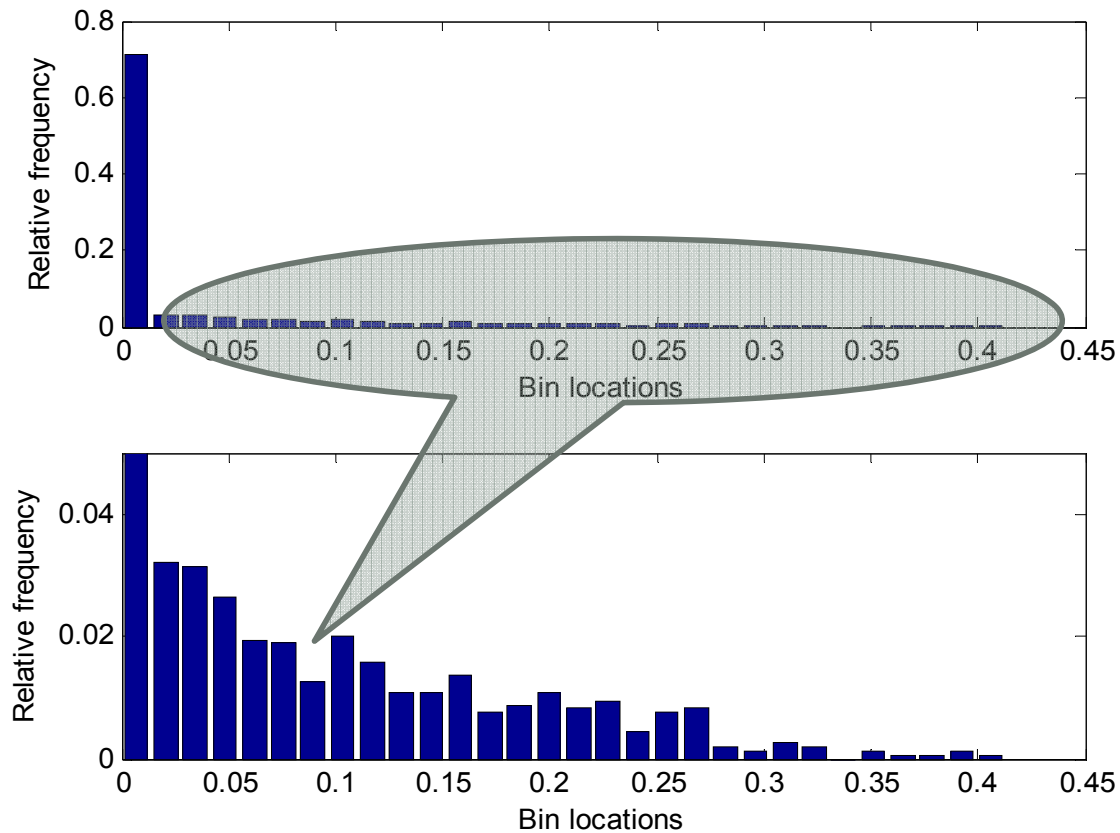
- A randomly generated serial line



- Accuracy evaluation:
 - 2000 lines generated
 - 80%: Single BN cases
 - 20%: Multiple BN cases. 70%: true BN machine is in candidate set.

Effectiveness

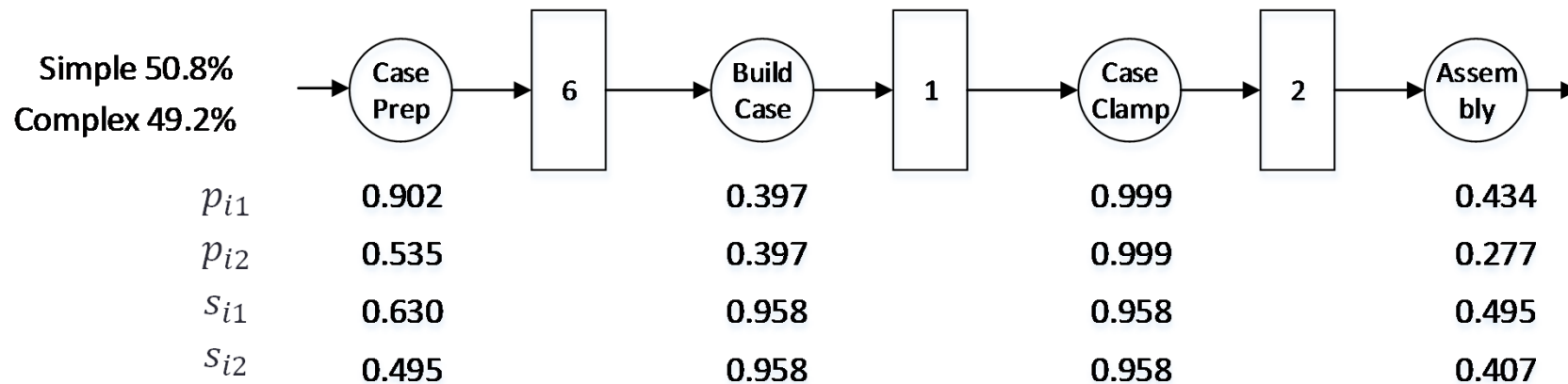
- Simulation comparison:
 - Single BN cases: True and identified BNs, difference of derivatives



- Multiple BN cases
 - Similar histogram pattern. Avg difference using severity index: 0.06

A Case Study

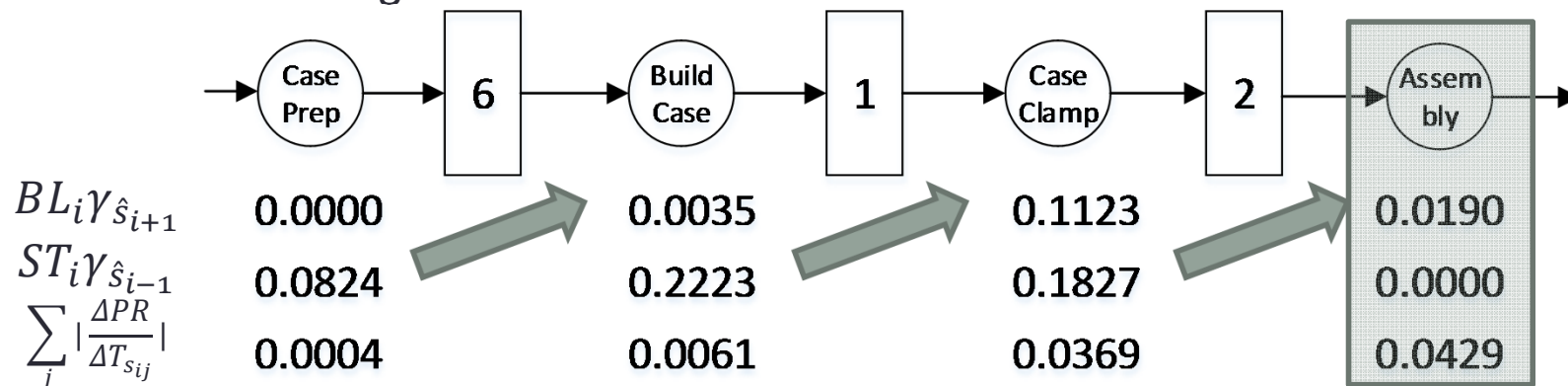
- Techline furniture assembly line.
 - Build main body in a serial line, assemble with drawers and doors at the end. Manual operations.
 - Feeding and shipping inventory large enough.
 - Simple and complex product groups.
 - Setup used for workers to read instructions and prepare tools.



- Throughput 205 (from production count) vs 218 (from model)
 - Model validated

Bottleneck Analysis

- Bottleneck indicators
 - Setup bottleneck product
 - $|0.5 - \alpha_1| = |0.5 - \alpha_2|$, same effectiveness, but reduce setup time for complex product is more practical.
 - Bottleneck Machine
 - Arrow assignment finds the true bottleneck machine



- Reduce 20% setup time on simple and complex assembly station
 - The throughput increase 3.21% (7 products).

Summary

- Effective tools to evaluate flexible production systems with setups are developed.
- Bottlenecks setup machine and bottleneck product are studied.
- Application in a furniture assembly line is carried out.

- Future work
 - Other reliability models (exponential, general)
 - Product sequencing and production control
 - Transient analysis

Call for Papers

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- **Guest Editors:**
 - Prof. Jingshan Li, University of Wisconsin-Madison
 - Prof. Chrissoleon Papadopoulos, Aristotle University of Thessaloniki
 - Prof. Liang Zhang, University of Connecticut
- **Important Dates:**
 - **Manuscript submission:** **June 30, 2015**
 - Tentative publication date: September 2016